

# BOUNDS ON SOME VAN DER WAERDEN NUMBERS

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## Abstract

For positive integers  $s$  and  $k_1, k_2, \dots, k_s$ , the van der Waerden number  $w(k_1, k_2, \dots, k_s; s)$  is the minimum integer  $n$  such that for every  $s$ -coloring of  $\{1, 2, \dots, n\}$ , with colors  $1, 2, \dots, s$ , there is a  $k_i$ -term arithmetic progression of color  $i$  for some  $i$ . We give an asymptotic lower bound for  $w(k, m; 2)$  for fixed  $m$ . We include a table of values of  $w(k, 3; 2)$  that are very close to this lower bound for  $m = 3$ . We also give a lower bound for  $w(k, k, \dots, k; s)$  that slightly improves previously-known bounds. Upper bounds for  $w(k, 4; 2)$  and  $w(4, 4, \dots, 4; s)$  are also provided.

## 1. Introduction

Two fundamental theorems in combinatorics are van der Waerden's Theorem [18] and Ramsey's Theorem [16]. The theorem of van der Waerden says that for all positive integers  $s$  and  $k_1, k_2, \dots, k_s$ , there exists a least positive integer  $n = w(k_1, k_2, \dots, k_s; s)$  such that whenever  $[1, n] = \{1, 2, \dots, n\}$  is  $s$ -colored (i.e., partitioned into  $s$  sets), there is a  $k_i$ -term arithmetic progression with color  $i$  for some  $i$ ,  $1 \leq i \leq s$ .

Similarly, Ramsey's Theorem has an associated "threshold" function  $R(k_1, k_2, \dots, k_s; s)$  (which we will not define here). The order of magnitude of  $R(k, 3; 2)$  is known to be  $\frac{k^2}{\log k}$  [11], while the best known upper bound on  $R(k, m; 2)$  is fairly close to the best known lower bound. In contrast, the order of magnitude of  $w(k, 3; 2)$  is not known, and the best known lower and upper bounds on  $w(k, k; 2)$  are

$$(k-1)2^{(k-1)} \leq w(k, k; 2) < 2^{2^{2^{2^{(k+9)}}}},$$

the lower bound known only when  $k - 1$  is prime. The lower bound is due to Berlekamp [1] and the upper bound is a celebrated result of Gowers [6], which answered a long-standing question of Ron Graham. Graham currently offers 1000 USD for a proof or disproof of  $w(k, k; 2) < 2^{k^2}$  [2]. Several other open problems are stated in [14].

Recently, there have been some further breakthroughs in the study of the van der Waerden function  $w(k, m; 2)$ . One was the amazing calculation that  $w(6, 6; 2) = 1132$  by Kouril [12], extending the list of previously known values  $w(3, 3; 2) = 9$ ,  $w(4, 4; 2) = 35$ , and  $w(5, 5; 2) = 178$ . A list of other known exact values of  $w(k, m; 2)$  appears in [15]. Improved lower bounds on several specific values of  $w(k, k; s)$  are given in [3] and [10].

In another direction, Graham [7] gives an elegant proof that if one defines  $w_1(k, 3)$  to be the least  $n$  such that every 2-coloring of  $[1, n]$  gives either  $k$  consecutive integers in the first color or a 3-term arithmetic progression in the second color, then

$$k^{c \log k} < w_1(k, 3) < k^{dk^2},$$

for suitable constants  $c, d > 0$ . This immediately gives  $w(k, 3; 2) < k^{dk^2}$  since we trivially have  $w(k, 3; 2) \leq w_1(k, 3)$ . In view of Graham's bounds on  $w_1(k, 3)$ , it would be desirable to obtain improved bounds on  $w(k, 3; 2)$ . Of particular interest is the question of whether or not there is a non-polynomial lower bound for  $w(k, 3; 2)$ .

In this note we give a lower bound of  $w(k, 3; 2) > k^{(2-o(1))}$ . Although this may seem weak, we do know that  $w(k, 3; 2) < k^2$  for  $5 \leq k \leq 16$  (i.e., for all known values of  $w(k, 3; 2)$  with  $k \geq 5$ ; see Table 1). More generally, we give a lower bound on  $w(k, m; 2)$  for arbitrary fixed  $m$ . We also present a lower bound for the classical van der Waerden numbers  $w(k, k, \dots, k; s)$  that is a slight improvement over previously published bounds. In addition, we present an upper bound for  $w(k, 4; 2)$  and an upper bound for  $w(4, 4, \dots, 4; s)$ .

## 2. Upper and Lower Bounds for Certain van der Waerden Functions

We shall need several definitions, which we collect here.

For positive integers  $k$  and  $n$ ,

$$r_k(n) = \max_{S \subseteq [1, n]} \{|S| : S \text{ contains no } k\text{-term arithmetic progression}\}.$$

For positive integers  $k$  and  $m$ , denote by  $\chi_k(m)$  the minimum number of colors required to color  $[1, m]$  so that there is no monochromatic  $k$ -term arithmetic progression.

The function  $w_1(k, 3)$  has been defined in Section 1. Similarly, we define  $w_1(k, 4)$  to be the least  $n$  such that every 2-coloring of  $[1, n]$  has either  $k$  consecutive integers in the first color or a 4-term arithmetic progression in the second color.

We begin with an upper bound for  $w_1(k, 4)$ . The proof is essentially the same as the proof given by Graham [7] of an upper bound for  $w_1(k, 3)$ . For completeness, we include the proof here. We will make use of a recent result of Green and Tao [9], who showed that for some constant  $c > 0$ ,

$$r_4(n) < ne^{-c\sqrt{\log \log n}} \tag{1}$$

for all  $n \geq 3$ .

**Proposition 2.1** There exists a constant  $c > 0$  such that  $w_1(k, 4) < e^{k^{c \log k}}$  for all  $k \geq 2$ .

*Proof.* Suppose we have a 2-coloring of  $[1, n]$  (assume  $n \geq 4$ ) with no 4-term arithmetic progression of the second color and no  $k$  consecutive integers of the first color. Let  $t_1 < t_2 < \dots < t_m$  be the integers of the second color. Hence,  $m < r_4(n)$ . Let us define  $t_0 = 0$  and  $t_{m+1} = n$ . Then there must be some  $i$ ,  $1 \leq i \leq m$ , such that

$$t_{i+1} - t_i > \frac{n}{2r_4(n)}.$$

(Otherwise, using  $r_4(n) \geq 3$ , we would have  $n = \sum_{i=0}^m (t_{i+1} - t_i) \leq \frac{n(m+1)}{2r_4(n)} \leq \frac{n(r_4(n)+1)}{2r_4(n)} \leq \frac{n}{2} + \frac{n}{6}$ .)

Using (1), we now have an  $i$  with

$$t_{i+1} - t_i > \frac{n}{2r_4(n)} > \frac{1}{2}e^{c\sqrt{\log \log n}}.$$

If  $n \geq e^{k^{d \log k}}$ ,  $d = c^{-2}$ , then  $\frac{1}{2}e^{c\sqrt{\log \log n}} \geq k$  and we have  $k$  consecutive integers of the first color, a contradiction. Hence,  $n < e^{k^{d \log k}}$  and we are done.  $\square$

Clearly  $w(k, 4; 2) \leq w_1(k, 4)$ . Consequently, we have the following result.

**Corollary 2.2** There exists a constant  $d > 0$  such that  $w(k, 4; 2) < e^{k^{d \log k}}$  for all  $k \geq 2$ .

Using Green and Tao's result, it is not difficult to obtain an upper bound for  $w(4, 4, \dots, 4; s)$ .

**Proposition 2.3** There exists a constant  $d > 0$  such that  $w(4, 4, \dots, 4; s) < e^{s^{d \log s}}$  for all  $s \geq 2$ .

*Proof.* Consider a  $\chi_4(m)$ -coloring of  $[1, m]$  for which there is no monochromatic 4-term arithmetic progression. Some color must be used at least  $\frac{m}{\chi_4(m)}$  times, and hence  $\frac{m}{\chi_4(m)} \leq r_4(m)$  so that  $\frac{m}{r_4(m)} \leq \chi_4(m)$ . Let  $c > 0$  be such that (1) holds for all  $n \geq 3$ , and let  $m = e^{s^{d \log s}}$ , where  $d = c^{-2}$ . Then  $\chi_4(m) \geq \frac{m}{r_4(m)} > e^{c\sqrt{\log \log m}} = s$ . This means that every  $s$ -coloring of  $[1, m]$  has a monochromatic 4-term arithmetic progression. Since  $m = e^{s^{d \log s}}$ , the proof is complete.  $\square$

It is interesting that the bounds in Corollary 2.2 and Proposition 2.3 have the same form.

The following theorem gives a lower bound on  $w(k, k, \dots, k; s)$ . It is deduced without too much difficulty from the Symmetric Hypergraph Theorem as it appears in [8], combined with an old result of Rankin [17]. To the best of our knowledge it has not appeared in print before, even though it is better, for large  $s$ , than the standard lower bound  $\frac{cs^k}{k}(1 + o(1))$  (see [8]), as well as the lower bounds  $s^{k+1} - \sqrt{c(k+1)\log(k+1)}$  and  $\frac{ks^k}{e^{(k+1)^2}}$  due to Erdős and Rado [4], and Everts [5], respectively. We give the proof in some detail. The proof makes use of the following facts:

$$\chi_k(n) < \frac{2n \log n}{r_k(n)}(1 + o(1)), \quad (2)$$

which appears in [8] as a consequence of the Symmetric Hypergraph Theorem; and

$$r_k(n) > ne^{-c(\log n)^{\frac{1}{\lceil \log_2 k \rceil + 1}}}, \quad (3)$$

which, for some constant  $c > 0$ , holds for all  $n \geq 3$  (this appears in [17]).

**Theorem 2.4** Let  $k \geq 3$  be fixed, and let  $z = \lfloor \log_2 k \rfloor$ . There exists a constant  $d > 0$  such that  $w(k, k, \dots, k; s) > s^{d(\log s)^z}$  for all sufficiently large  $s$ .

*Proof.* We make use of the observation that for positive integers  $s$  and  $m$ , if  $s \geq \chi_k(m)$ , then  $w(k, k, \dots, k; s) > m$ , which is clear from the definitions. For large enough  $m$ , (2) gives

$$\chi_k(m) < \frac{2m \log m}{r_k(m)} \left(1 + \frac{1}{2}\right) = \frac{3m \log m}{r_k(m)}. \quad (4)$$

Now let  $d = \left(\frac{1}{2c}\right)^{z+1}$ , where  $c$  is from (3), and let  $m = s^{d(\log s)^z}$ , where  $s$  is large enough so that (4) holds. By (3), noting that  $\log m = d(\log s)^{z+1} = \left(\frac{\log s}{2c}\right)^{z+1}$ , we have

$$\frac{m}{r_k(m)} < e^{c(\log m)^{\frac{1}{z+1}}} = e^{c \cdot \frac{\log s}{2c}} = \sqrt{s}.$$

Therefore,

$$\frac{3m \log m}{r_k(m)} < 3d\sqrt{s}(\log s)^{z+1} < s$$

for sufficiently large  $s$ . Thus, for sufficiently large  $s$ ,

$$\chi_k(m) < \frac{3m \log m}{r_k(m)} < s.$$

According to the observation at the beginning of the proof, this implies that  $w(k, k, \dots, k; s) > m = s^{d(\log s)^z}$ , as required.  $\square$

We now give a lower bound on  $w(k, m; 2)$ . We make use of the Lovász Local Lemma (see [8] for a proof), which will be implicitly stated in the proof.

**Theorem 2.5** Let  $m \geq 3$  be fixed. Then for all sufficiently large  $k$ ,

$$w(k, m; 2) > k^{m-1-\frac{1}{\log \log k}}.$$

*Proof.* Given  $m$ , choose  $k > m$  large enough so that

$$k^{\frac{1}{2m \log \log k}} > \left( m - \frac{1}{2 \log \log k} \right) \log k \tag{5}$$

and

$$6 < \frac{\log k}{\log \log k}. \tag{6}$$

Next, let  $n = \lfloor k^{m-1-\frac{1}{\log \log k}} \rfloor$ . To prove the theorem, we will show that there exists a (red, blue)-coloring of  $[1, n]$  for which there is no red  $k$ -term arithmetic progression and no blue  $m$ -term arithmetic progression.

For the purpose of using the Lovász Local Lemma, randomly color  $[1, n]$  in the following way. For each  $i \in [1, n]$ , color  $i$  red with probability  $p = 1 - k^{\alpha-1}$  where

$$\alpha \stackrel{\text{def}}{=} \frac{1}{2m \log \log k},$$

and color it blue with probability  $1 - p$ .

Let  $\mathcal{P}$  be any  $k$ -term arithmetic progression. Then, since  $1 + x \leq e^x$ , the probability that  $\mathcal{P}$  is red is

$$p^k = (1 - k^{\alpha-1})^k \leq (e^{-k^{\alpha-1}})^k = e^{-k^\alpha}.$$

Hence, applying (5), we have

$$p^k < \left( \frac{1}{e} \right)^{\left( m - \frac{1}{2 \log \log k} \right) \log k} = \frac{1}{k^{m - \frac{1}{2 \log \log k}}}.$$

Also, for any  $m$ -term arithmetic progression  $\mathcal{Q}$ , the probability that  $\mathcal{Q}$  is blue is

$$(1 - p)^m = (k^{\alpha-1})^m = \frac{1}{k^{m - \frac{1}{2 \log \log k}}}.$$

Now let  $\mathcal{P}_1, \mathcal{P}_2, \dots, \mathcal{P}_t$  be all of the arithmetic progressions in  $[1, n]$  with length  $k$  or  $m$ . So that we may apply the Lovász Local Lemma, we form the “dependency graph”  $G$  by setting  $V(G) = \{\mathcal{P}_1, \mathcal{P}_2, \dots, \mathcal{P}_t\}$  and  $E(G) = \{\{\mathcal{P}_i, \mathcal{P}_j\} : i \neq j, \mathcal{P}_i \cap \mathcal{P}_j \neq \emptyset\}$ . For each  $\mathcal{P}_i \in V(G)$ , let  $d(\mathcal{P}_i)$  denote the degree of the vertex  $\mathcal{P}_i$  in  $G$ , i.e.,  $|\{e \in E(G) : \mathcal{P}_i \in e\}|$ . We now estimate  $d(\mathcal{P}_i)$  from above. Let  $x \in [1, n]$ . The number of  $k$ -term arithmetic progressions  $\mathcal{P}$  in  $[1, n]$  that contain  $x$  is bounded above by  $k \cdot \frac{n}{k-1}$ , since there are  $k$  positions that  $x$  may occupy in  $\mathcal{P}$  and since the gap size of  $\mathcal{P}$  cannot exceed  $\frac{n}{k-1}$ . Similarly, the number of  $m$ -term arithmetic progressions  $\mathcal{Q}$  in  $[1, n]$  that contain  $x$  is bounded above by  $m \cdot \frac{n}{m-1}$ .

Let  $\mathcal{P}_i$  be any  $k$ -term arithmetic progression contained in  $[1, n]$ . The total number of  $k$ -term arithmetic progressions  $\mathcal{P}$  and  $m$ -term arithmetic progressions  $\mathcal{Q}$  in  $[1, n]$  that may have non-empty intersection with  $\mathcal{P}_i$  is bounded above by

$$k \left( k \cdot \frac{n}{k-1} + m \cdot \frac{n}{m-1} \right) < kn \left( 2 + \frac{2}{m-1} \right), \quad (7)$$

since  $k > m$ . Thus,  $d(\mathcal{P}_i) < kn \left( 2 + \frac{2}{m-1} \right)$  when  $|\mathcal{P}_i| = k$ . Likewise,  $d(\mathcal{P}_i) < mn \left( 2 + \frac{2}{m-1} \right)$  when  $|\mathcal{P}_i| = m$ . Thus, for all vertices  $\mathcal{P}_i$  of  $G$ , we have  $d(\mathcal{P}_i) < kn \left( 2 + \frac{2}{m-1} \right)$ .

To finish setting up the hypotheses for the Lovász Local Lemma, we let  $X_i$  denote the event that the arithmetic progression  $\mathcal{P}_i$  is red if  $|\mathcal{P}_i| = k$ , or blue if  $|\mathcal{P}_i| = m$ , and we let  $d = \max_{1 \leq i \leq t} d(\mathcal{P}_i)$ . We have seen above that for all  $i$ ,  $1 \leq i \leq t$ , the probability that  $X_i$  occurs is at most

$$q \stackrel{\text{def}}{=} \frac{1}{k^{m - \frac{1}{2 \log \log k}}},$$

while from (7) we have  $d < 2kn \left( 1 + \frac{1}{m-1} \right)$ .

We are now ready to apply the Lovász Local Lemma, which says that in these circumstances, if the condition  $eq(d+1) < 1$  is satisfied, then there is a (red, blue)-coloring of  $[1, n]$  such that no event  $X_i$  occurs, i.e., such that there is no red  $k$ -term arithmetic progression and no blue  $m$ -term arithmetic progression. This will imply

$$w(k, m; 2) > n = k^{m-1 - \frac{1}{\log \log k}},$$

as desired. Thus, the proof will be complete when we verify that  $eq(d+1) < 1$ . Using  $m \geq 3$ , we have  $d < 3kn$ , so that  $d+1 < 3kn+1 < e^2 kn$ . Hence, it is sufficient to verify that

$$e^3 qkn < 1. \quad (8)$$

Since  $q = \frac{1}{k^{m - \frac{1}{2 \log \log k}}}$  and  $n \leq k^{m-1 - \frac{1}{\log \log k}}$ , inequality (8) may be reduced to (6), and the proof is now complete.  $\square$

*Remark.* As long as  $k > e^{m^6}$ , the inequality of Theorem 2.5 holds. To show this, we need only show that conditions (5) and (6) hold if  $k > e^{m^6}$ . That (6) holds is obvious. For (5), it suffices to have  $k^{\frac{1}{2m \log \log k}} > m \log k$ ; that is  $\log k > 2m \log \log k (\log m + \log \log k)$ . When  $k \geq e^{m^6}$ , we have  $2m \log \log k (\log m + \log \log k) \leq 2(\log k)^{1/6} \log \log k \left( \frac{1}{6} \log \log k + \log \log k \right) = \frac{7}{3} (\log k)^{1/6} (\log \log k)^2$ . Since  $(\log \log k) < (\log k)^{7/20}$  for  $k \geq e^{m^6}$  we have  $2m \log \log k (\log m + \log \log k) \leq \frac{7}{3} (\log k)^{13/15}$ . Finally, since  $(\log k)^{2/15} \geq \frac{7}{3}$  for  $k \geq e^{m^6}$ , condition (5) is satisfied.

We end with a table of computed values. These were all computed with a standard backtrack algorithm except for  $w(14, 3; 2)$ ,  $w(15, 3; 2)$ , and  $w(16, 3; 2)$ , which are due to Michal Kouril [13]. The values  $w(k, 3; 2)$ ,  $k \leq 12$ , appeared previously in [15].

$k$	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
$w(k, 3; 2)$	6	9	18	22	32	46	58	77	97	114	135	160	186	218	238
$w_1(k, 3)$	9	23	34	73	113	193	?	?	?	?	?	?	?	?	?

**Table 1: Small values of  $w(k, 3)$  and  $w_1(k, 3)$**

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